

C. Amendments to the Claims

Claim 1 (currently amended): A method for generating prime numbers for use in encryption, the method comprising the steps of:

- a. generating a random number 'n';
- b. checking if the random number 'n' is an exact power of another positive integer;
if the random number 'n' is an exact power of the another positive integer, then:
- c. declaring the random number 'n' to be composite; and
if the random number 'n' is not an exact power of the another positive integer, then:
- d. performing an extension ring test on the random number 'n'.

Claim 2 (original): The method as recited in claim 1 wherein the step of performing the extension ring test comprises the steps of:

- a. choosing a set of polynomials $g(x)$;
- b. choosing a polynomial $f(x)$;
if $[g(x)]^n \neq g(x^n) \bmod(f(x), n)$, for the chosen $f(x)$ and any $g(x)$ belonging to the chosen set of polynomials $g(x)$ then:
- c. declaring the random number 'n' to be composite; and
if $[g(x)]^n = g(x^n) \bmod(f(x), n)$, for the chosen $f(x)$ and all $g(x)$ belonging to the chosen set of polynomials $g(x)$ then:
- d. declaring the random number 'n' to be prime.

Claim 3 (original): The method as recited in claim 1 wherein the step of performing the extension ring test comprises the steps of:

a. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

- i. the number 'r' is a prime number;
- ii. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;
- iii. $n^{(r-1)/q} \not\equiv 1 \pmod{r}$; and
- iv. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step a exists, then for all values of the number 'r' less than the random number 'n':

b. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1;

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

- i. declaring the number 'n' to be composite;

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

- ii. declaring the number 'n' to be prime; and

if a number 'r' satisfying the conditions specified in step a exists, then:

c. checking whether $(x + a)^n \equiv x^n + a \pmod{(n, x^r - 1)}$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$;

if $(x + a)^n \not\equiv x^n + a \pmod{(n, x^r - 1)}$ for any integer value of 'a' between 1 and $(2\sqrt{r} \log_2 n)$, then:

- i. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \pmod{(n, x^r - 1)}$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$, then:

ii. declaring the random number 'n' to be prime.

Claim 4 (currently amended): A method for generating prime numbers for use in encryption, the method comprising the steps of:

a. generating a random number 'n';

b. checking if the random number 'n' is an exact power of another positive integer;

if the random number 'n' is an exact power of the another positive integer, then:

c. declaring the random number 'n' to be composite; and

if the random number 'n' is not an exact power of the another positive integer, then:

d. performing an extension ring test, the extension ring test comprising:

i. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

1. the number 'r' is a prime number;

2. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;

3. $n^{(r-1)/q} \not\equiv 1 \pmod{r}$; and

4. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step a i exists, then

for all values of the number 'r' less than the random number 'n':

ii. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

1. declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

2. declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step a i exists, then:

iii. checking whether $(x + a)^n = x^n + a \bmod (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2^{\sqrt{r}} \log_2 n)$;

if $(x + a)^n \neq x^n + a \bmod (n, x^r - 1)$ for any integer value of 'a' between 1 and $(2^{\sqrt{r}} \log_2 n)$, then:

1. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \bmod (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2^{\sqrt{r}} \log_2 n)$, then:

2. declaring the random number 'n' to be prime.

Claim 5 (currently amended): The method as recited in claim 1, wherein the method is implemented as a computer program product.

Claim 6 (canceled)

Claim 7 (Currently amended): A method of for deterministically testing primality of a random number 'n' in polynomial time for use in encryption, the method comprising the steps of:

a. checking if the random number 'n' is an exact power of another positive integer;

if the random number 'n' is an exact power of the another positive integer, then:

b. declaring the random number 'n' to be composite; and

if the random number 'n' is not an exact power of the another positive integer, then:

c. performing an extension ring test.

Claim 8 (original): The method as recited in claim 7 wherein the step of performing the extension ring test comprises the steps of:

a. choosing a set of polynomials $g(x)$;

b. choosing a polynomial $f(x)$;

if $[g(x)]^n \neq g(x^n) \bmod(f(x), n)$, for the chosen $f(x)$ and any $g(x)$ belonging to the chosen set of polynomials $g(x)$ then:

c. declaring the random number 'n' to be composite; and

if $[g(x)]^n = g(x^n) \bmod(f(x), n)$, for the chosen $f(x)$ and all $g(x)$ belonging to the chosen set of polynomials $g(x)$ then:

d. declaring the random number 'n' to be prime.

Claim 9 (original): The method as recited in claim 7 wherein the step of performing the extension ring test comprises the steps of:

a. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

i. the number 'r' is a prime number;

ii. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;

iii. $n^{(r-1)/q} \neq 1 \bmod(r)$; and

iv. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step a exists, then for all values of the number 'r' less than the random number 'n':

- b. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

- i. declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

- ii. declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step a exists, then:

- c. checking whether $(x + a)^n = x^n + a \text{ mod } (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$;

if $(x + a)^n \neq x^n + a \text{ mod } (n, x^r - 1)$ for any integer value of 'a' between 1 and $(2\sqrt{r} \log_2 n)$, then:

- i. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \text{ mod } (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$, then:

- ii. declaring the random number 'n' to be prime.

Claim 10 (currently amended): A method of for deterministically testing primality of a random number 'n' in polynomial time for use in encryption, the method comprising the steps of:

- a. checking if the random number 'n' is an exact power of another positive integer;

if the random number 'n' is an exact power of the another positive integer, then:

b. declaring the random number 'n' to be composite; and

if the random number 'n' is not an exact power of the another positive integer, then:

c. performing an extension ring test, the extension ring test comprising:

i. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

1. the number 'r' is a prime number;
2. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;
3. $n^{(r-1)/q} \not\equiv 1 \pmod{r}$; and
4. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step a i exists, then
for all values of the number 'r' less than the random number 'n':

ii. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

1. declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

2. declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step a i exists, then:

iii. checking whether $(x + a)^n = x^n + a \bmod (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$;

if $(x + a)^n \neq x^n + a \bmod (n, x^r - 1)$ for any integer value of 'a' between 1 and $(2\sqrt{r} \log_2 n)$, then:

1. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \bmod (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$, then:

2. declaring the random number 'n' to be prime.

Claim 11 (currently amended): The method as recited in claim 7, wherein the method is implemented as a computer program product.

Claim 12 (canceled)

Claim 13 (currently amended): A system for deterministically testing primality of a random number 'n' in polynomial time for use in encryption, the system comprising: a primality tester, the primality tester performing an extension ring test on the random number 'n' so as to determine whether the random number 'n' is prime, wherein the extension ring test comprises:

a. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

i. the number 'r' is a prime number;

ii. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;

iii. $n^{(r-1)/q} \neq 1 \bmod (r)$; and

iv. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step a exists, then for all values of the number 'r' less than the random number 'n':

b. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

i. declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

ii. declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step a exists, then:

c. checking whether $(x + a)^n = x^n + a \bmod (n, x' - 1)$ for all integer values of 'a' from 1 to $(2^{\sqrt{r}} \log_2 n)$;

if $(x + a)^n \neq x^n + a \bmod (n, x' - 1)$ for any integer value of 'a' between 1 and $(2^{\sqrt{r}} \log_2 n)$, then:

i. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \bmod (n, x' - 1)$ for all integer values of 'a' from 1 to $(2^{\sqrt{r}} \log_2 n)$, then:

ii. declaring the random number 'n' to be prime.

~~a. means for checking if the random number 'n' is an exact power of another positive integer;~~

~~if the random number 'n' is an exact power of another positive integer, then:~~

~~b. means for declaring the random number 'n' to be composite; and~~

~~if the random number 'n' is not an exact power of another positive integer, and;~~

~~e. means for performing an extension ring test.~~

Claim 14 (canceled)

Claim 15 (canceled)

Claim 16 (currently amended): A system for encrypting a communication prime number generator, the prime number generator being connected to an encryption system, the encryption system using prime numbers generated by the prime number generator in encryption, said system including a the prime number generator, wherein the improvement comprises comprising:

~~a. the prime number generator having a random number generator, for the random number generator generating a random number 'n'; and~~

~~b. : means for checking if the random number 'n' is an exact power of another positive integer; if the random number 'n' is an exact power of another positive integer, then declaring the random number 'n' to be composite; if the random number 'n' is not an exact power of another positive integer, then a primality tester, the primality tester performing an extension ring test on the random number 'n' so as to determine whether the random number 'n' is prime.~~

Claim 17 (canceled)

Claim 18 (currently amended): The encryption system prime number generator as claimed in claim 16, wherein the extension ring test primality tester comprises means for:

~~a. generating a random number 'n';~~

~~b. checking if the random number 'n' is an exact power of another positive integer;~~

~~if the random number 'n' is an exact power of another positive integer, then:~~

~~c. declaring the random number 'n' to be composite; and~~

~~if the random number 'n' is not an exact power of another positive integer, then:~~

~~d. performing an extension ring test, the extension ring test comprising means for:~~

i. a. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

1. i. the number 'r' is a prime number;

2. ii. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;

3. iii. $n^{(r-1)/q} \not\equiv 1 \pmod{r}$; and

4. iv. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step a exists, then for all values of the number 'r' less than the random number 'n':

ii. b. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

1. i. means for declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

2. ii. means for declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step a exists, then:

iii. c. checking whether $(x + a)^n = x^n + a \text{ mod } (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$;

if $(x + a)^n \neq x^n + a \text{ mod } (n, x^r - 1)$ for any integer value of 'a' between 1 and $(2\sqrt{r} \log_2 n)$, then:

4. i. means for declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \text{ mod } (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$, then:

2. ii. means for declaring the random number 'n' to be prime.

Claim 19 (canceled)

Claim 20 (canceled)

Claim 21 (canceled)

Claim 22 (new): A computer program product for use with a computer, the computer program product comprising a computer usable medium having a computer readable program code embodied therein for generating prime numbers for use in encryption, the computer readable program code comprising instructions for:

a. generating a random number 'n';

b. checking if the random number 'n' is an exact power of another positive integer;

if the random number 'n' is an exact power of the another positive integer, then:

c. declaring the random number 'n' to be composite; and

if the random number 'n' is not an exact power of the another positive integer, then:

d. performing an extension ring test, the extension ring test comprising:

i. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

1. the number 'r' is a prime number;
2. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;
3. $n^{(r-1)/q} \not\equiv 1 \pmod{r}$; and
4. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step i exists, then for all values of the number 'r' less than the random number 'n':

ii. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

1. declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

2. declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step i exists, then:

iii. checking whether $(x + a)^n \equiv x^n + a \pmod{(n, x^r - 1)}$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$;

if $(x + a)^n \not\equiv x^n + a \pmod{(n, x^r - 1)}$ for any integer value of 'a' between 1 and $(2\sqrt{r} \log_2 n)$, then:

1. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \pmod{(n, x^r - 1)}$ for all integer values of 'a' from 1 to $(2\sqrt{r} \log_2 n)$, then:

2. declaring the random number 'n' to be prime.

Claim 23 (new): A computer program product for use with a computer, the computer program product comprising a computer usable medium having a computer readable program code embodied therein for deterministically testing primality of a random number 'n' in polynomial time for use in encryption, the computer readable program code comprising instructions for:

a. checking if the random number 'n' is an exact power of another positive integer;

if the random number 'n' is an exact power of the another positive integer, then:

b. declaring the random number 'n' to be composite; and

if the random number 'n' is not an exact power of the another positive integer, then:

c. performing an extension ring test, the extension ring test comprising:

i. determining the smallest number 'r' less than the random number 'n', the number 'r' satisfying the following conditions:

1. the number 'r' is a prime number;

2. the largest prime factor 'q' of $(r - 1)$ is greater than or equal to $(4\sqrt{r} \log_2 n)$;

3. $n^{(r-1)/q} \not\equiv 1 \pmod{r}$; and

4. the greatest common divisor of the number 'r' and the random number 'n' is equal to 1;

if no number 'r' satisfying all the conditions specified in step i exists, then for all values of the number 'r' less than the random number 'n':

ii. performing a check whether the greatest common divisor of the number 'r' and the random number 'n' is greater than 1; and

if the greatest common divisor of 'r' and 'n' is greater than 1 for any value of the number 'r', then:

1. declaring the number 'n' to be composite; and

if the greatest common divisor of 'r' and 'n' is equal to 1 for all values of the number 'r', then:

2. declaring the number 'n' to be prime;

if a number 'r' satisfying the conditions specified in step i exists, then:

iii. checking whether $(x + a)^n = x^n + a \bmod (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2^{\sqrt{r}} \log_2 n)$;

if $(x + a)^n \neq x^n + a \bmod (n, x^r - 1)$ for any integer value of 'a' between 1 and $(2^{\sqrt{r}} \log_2 n)$, then:

1. declaring the number 'n' to be composite; and

if $(x + a)^n = x^n + a \bmod (n, x^r - 1)$ for all integer values of 'a' from 1 to $(2^{\sqrt{r}} \log_2 n)$, then:

2. declaring the random number 'n' to be prime.